

## 119

state according to a Markov process, but takes an arbitrarily distributed amount of time between state changes. The further development of the process depends not only on the present state, but also on the amount of time that elapsed since entry into the state. The generalized Markov Renewal Process allows that some of the MRP states may not constitute regeneration points. Nakagawa and Osaki used it in modeling and analyzing reliable systems, and showed that it is an extremely powerful analyzing tool. Markovian analysis, in principle, could be performed directly without the intervention of SPN. In practice, however, it is difficult to determine the correct Markov model for even moderately complex systems without such a descriptive 'language' like ESPN. Using ESPN, the modeling flexibility of Petri Net and the analyzing power of Markov Renewal Process can be exploited more easily. Moreover, in this paper, we introduce an Abstract Partial Reachability Graph (APRG), a sub-set of Reachability Graph, to simplify the Markov solution.

## 2. PETRI NET AND STOCHASTIC PETRI NET

A Petri Net is an abstract, formal graph model useful for modeling systems that exhibit concurrent, asynchronous or non-deterministic behavior. Firstly, recall the composition of a Petri Net (PN) bipartite graph [1]: a set of *places*,  $P$  (drawn as circles) representing conditions, a set of *transitions*,  $T$  (drawn as bars) representing events, and a set of *arcs*,  $A$ , which connect places to transitions or transitions to places. Places may contain *tokens* (drawn as small filled circles) denoting the conditions holding at any given time. The state of a Petri Net, called the PN *marking*, is defined by the number of tokens contained in each place.

A place is an input to a transition if an arc exists from the place to transition; a place is an output to a transition if an arc exists from a transition to the place. A transition is *enabled* when each of its input places contains at least one token. Enabled transitions can *fire*, by removing one token from each input place and depositing one token in each output place. As an event is usually enabled by a combination of conditions, a transition is enabled by a combination of tokens in places. Arcs are used to signify which combination of conditions must hold for the event to occur and which combination of conditions holds after the event occurs. Thus the firing of a transition causes a change of state (produces a different marking) for the Petri Net.

Stochastic Petri Net (SPN) [2-5] is defined by associating an exponentially distributed firing time with each transition. Once a transition is enabled, an exponentially distributed amount of time elapses. If the transition is still enabled, it will then fire. An SPN can be analyzed by considering all possible markings (enumerations of the tokens in each place) and solving the resulting reachability graph as a Markov chain. The analysis of SPN model provides information about the system it represents, provided the model is a valid representation of the system under study, and the solution of the model is correct.

## 3. EXTENDED STOCHASTIC PETRI NET MODEL FOR A MULTI-ROBOT SYSTEM

In this paper, as a case study, we consider a system of four robots with parallel and cooperative motions. Robots 1, 2 and 3 operate parallelly and independently. They also operate on a buffer, but at one time only one buffer operation is permitted. In addition, after Robot 3 operating on the buffer, it has to do cooperative motion with Robot 4, which is always ready (standby) to do the cooperative motion.

Figure 1 is the multi-robot system represented by Petri Net, which is referred to as an Extended Stochastic Petri Net in this paper because we allow the incorporation of arbitrary firing time distribution into it. Meanings of the places and transitions are shown as follows.

- $\varepsilon_1$ : Robot 1 is doing independent operation.
- $\varepsilon_2$ : Robot 1 is doing buffer operation.
- $\varepsilon_3$ : Robot 2 is doing independent operation.
- $\varepsilon_4$ : Robot 2 is doing buffer operation.
- $\varepsilon_5$ : Robot 3 is doing independent operation.
- $\varepsilon_6$ : Robot 3 is doing buffer operation.
- $\varepsilon_7$ : Robots 3 and 4 are doing cooperative motion.

- $\varepsilon_8$ : Robot 4 is standby.  
 $\varepsilon_9$ : Buffer is idle.  
 $\varepsilon_{10}$ : Buffer is busy.  
 $\tau_1$ : Robot 1 finishes independent operation and begins buffer operation.  
 $\tau_2$ : Robot 1 finishes buffer operation and begins independent operation once more.  
 $\tau_3$ : Robot 2 finishes independent operation and begins buffer operation.  
 $\tau_4$ : Robot 2 finishes buffer operation and begins independent operation once more.  
 $\tau_5$ : Robot 3 finishes independent operation and begins buffer operation.  
 $\tau_6$ : Robot 3 finishes buffer operation. Robots 3 and 4 begin cooperative motion.  
 $\tau_7$ : Robots 3 and 4 finish cooperative motion. Robot 3 begins independent operation once more and Robot 4 returns to standby status.

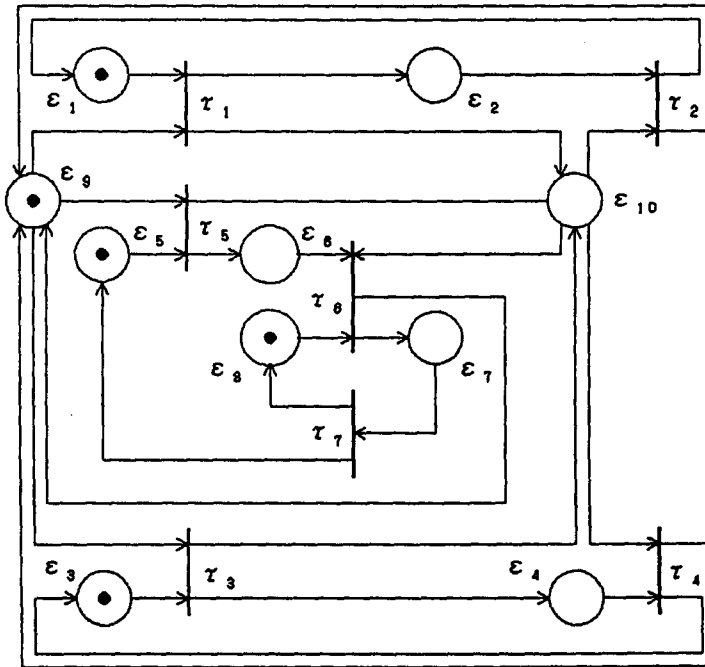


Figure 1. ESPN model of the multi-robot system.

Since disturbance may occur during Robots 3 and 4 doing their cooperative motion (i.e., the operation of transition  $\tau_7$ ) in this ESPN model of the multi-robot system, consider that a non-exponential distribution is associated with the firing time distribution of  $\tau_7$ . For simplification of numerical solution, all the other transitions  $\tau_i (i = 1, 2, \dots, 6)$  are assumed to obey exponential firing time distribution  $F_i(t) = 1 - e^{-\lambda_i t} (i = 1, 2, \dots, 6)$ .

An SPN model is generally analyzed by considering all possible marking (enumerations of the tokens in each place) and solving the resulting Reachability Graph (RG) or the marking process as a Markov chain. In this paper, for no longer all the firing times of transitions in the ESPN model obey exponential distribution, we define the marking process in terms of a renewal state space of a Markov Renewal Process, and obtain the analytic solution in the context of a generalized MRP [9]. Figure 2 is the Reachability Graph for the ESPN model in Figure 1, where the elements of the vector: 1 means a token in the corresponding place and 0 means no token. Each state  $S_i (i = 0, 1, \dots, 6)$  has the definition as follows.

- $S_0$ : Robots 1, 2 and 3 are on independent operation, Robot 4 is standby, and Buffer is idle.  
 $S_1$ : Robot 1 is on buffer operation, Robots 2 and 3 are on independent operation, Robot 4 is standby, and Buffer is busy.  
 $S_2$ : Robots 1 and 3 are on independent operation, Robot 2 is on buffer operation, Robot 4 is standby, and Buffer is busy.

- $S_3$ : Robots 1 and 2 are on independent operation, Robot 3 is on buffer operation, Robot 4 is standby, and Buffer is busy.
- $S_4$ : Robots 1 and 2 are on independent operation, Robots 3 and 4 are on cooperative motion, and Buffer is idle.
- $S_5$ : Robot 1 is on buffer operation, Robot 2 is on independent operation, Robots 3 and 4 are on cooperative motion, and Buffer is busy.
- $S_6$ : Robot 1 is on independent operation, Robot 2 is on buffer operation, Robots 3 and 4 are on cooperative motion, and Buffer is busy.

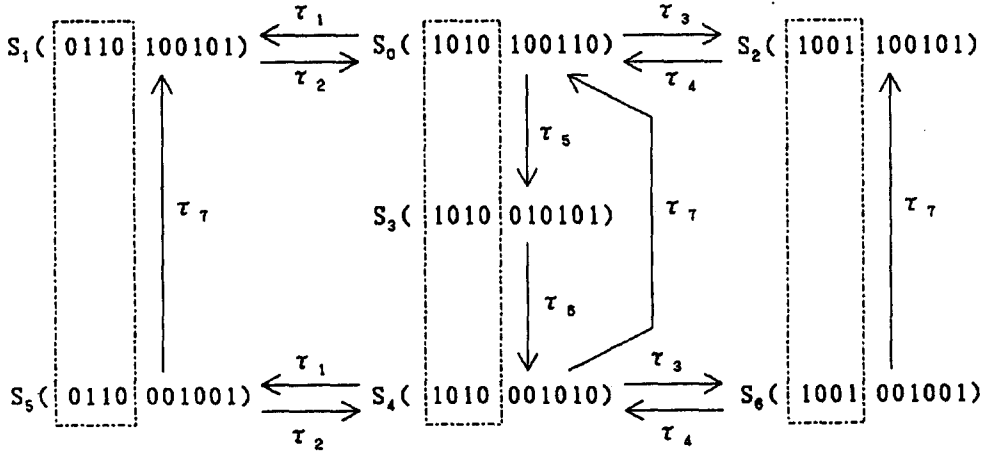


Figure 2. Reachability graph for the ESPN model in Figure 1.

#### 4. ANALYSIS OF ESPN MODEL USING A GENERALIZED MRP

Since non-exponential distribution is incorporated to the firing time of transition  $\tau_7$  in the ESPN model, the Markovian analytic solution is different from the solution by the Markov chain and becomes more complex to solve. We introduce the Abstract Partial Reachability Graph (APRG) to simplify the Markov solution.

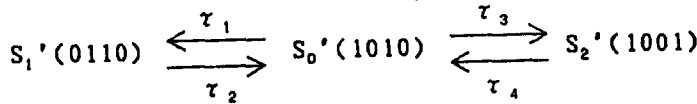


Figure 3. APRG derived from Figure 2.

First, let us just consider the operations of Robots 1 and 2. From Figure 2, we can see  $S_0$ ,  $S_3$  and  $S_4$  define the same state of Robots 1 and 2 on independent operations. We make them into a new state  $S'_0$ . Similarly, we find  $S_1$  and  $S_5$  defining the same state of Robot 1 on buffer operation and Robot 2 on independent operation, making them into a new state  $S'_1$ ,  $S_2$  and  $S_6$  defining the same state of Robot 1 on independent operation and Robot 2 on buffer operation, making them into a new state  $S'_2$ . The Reachability Graph of decreased states is re-drawn in Figure 3. It is called Abstract Partial Reachability Graph, which is equivalent to a simple Markov Renewal Process [9]. With the APRG in Figure 3, we need only to investigate the state transfers with all transitions ( $\tau_1$ ,  $\tau_2$ ,  $\tau_3$  and  $\tau_4$ ) of exponential firing time distributions.

For the APRG in Figure 3, transference probabilities between states  $S'_0$  and  $S'_j$  ( $j = 0, 1, 2$ ) are

$$p'_{00}(s) = \frac{1 - q'_{01}(s) - q'_{02}(s)}{1 - g'_{00}(s)} \quad (1)$$

$$p'_{0j}(s) = \frac{q'_{0j}(s)[1 - q'_{j0}(s)]}{1 - g'_{00}(s)}, \quad (j = 1, 2) \quad (2)$$

where,  $p'_{0j}(s)$  ( $j = 0, 1, 2$ ) are the Laplace-Stieljes (LS) Transforms of  $P'_{0j}(t)$ , and  $P'_{00}(t)$  means the probability of recurrent state transfer from  $S'_0$  to  $S'_0$ ;

$$g'_{00}(s) = q'_{01}(s)q'_{10}(s) + q'_{02}(s)q'_{20}(s) \quad (3)$$

is the recurrent first-passage time distribution of state  $S'_0$ ;

$$q'_{01}(s) = \frac{\lambda_1}{s + \lambda_1 + \lambda_3} \quad (4)$$

$$q'_{02}(s) = \frac{\lambda_3}{s + \lambda_1 + \lambda_3} \quad (5)$$

$$q'_{10}(s) = q'_{20}(s) = \frac{\lambda_b}{s + \lambda_b} \quad (6)$$

are the mass functions and here we have assumed

$$\lambda_2 = \lambda_4 = \lambda_b \quad (7)$$

Therefore, we have

$$p'_{00}(s) = \frac{s + \lambda_b}{s + \lambda_1 + \lambda_3 + \lambda_b} \quad (8)$$

$$p'_{01}(s) = \frac{\lambda_1}{s + \lambda_1 + \lambda_3 + \lambda_b} \quad (9)$$

$$p'_{02}(s) = \frac{\lambda_3}{s + \lambda_1 + \lambda_3 + \lambda_b} \quad (10)$$

Then let us go back to investigate the Reachability Graph of the ESPN model in Figure 2. From Figure 2 we see the RG forms a generalized Markov Renewal Process [9]. For we have assumed  $F_i(t)$  ( $i = 1, 2, \dots, 6$ ) are exponential distributions, states  $S_0, S_1, S_2, S_3$  and  $S_4$  are regeneration points. On the other hand, we assume that  $F_7(t)$  obeys Gamma distribution (shape parameter = 2)

$$F_7(t) = 1 - (1 + \lambda_7 t)e^{-\lambda_7 t} \quad (11)$$

Note that the firing distribution of transition  $\tau_7$ , which means Robots 3 and 4 finish cooperative motion, starts from state  $S_4$  not state  $S_5$  or  $S_6$ . We cannot determine the start point of states  $S_5$  and  $S_6$ . Therefore, we treat states  $S_5$  and  $S_6$  as non-regeneration points.

Thus, we obtain the mass functions of one-step state transference as follows

$$q_{01}(s) = \frac{\lambda_1}{s + \lambda_1 + \lambda_3 + \lambda_5} \quad (12)$$

$$q_{02}(s) = \frac{\lambda_3}{s + \lambda_1 + \lambda_3 + \lambda_5} \quad (13)$$

$$q_{03}(s) = \frac{\lambda_5}{s + \lambda_1 + \lambda_3 + \lambda_5} \quad (14)$$

$$q_{10}(s) = q_{20}(s) = q_{34}(s) = \frac{\lambda_b}{s + \lambda_b} \quad (15)$$

$$q_{40}(s) = \frac{\lambda_7^2}{\lambda_1 + \lambda_3 + \lambda_b} \left\{ \frac{\lambda_b}{(s + \lambda_7)^2} + \frac{\lambda_1 + \lambda_3}{(s + \lambda_1 + \lambda_3 + \lambda_5 + \lambda_7)^2} \right\} \quad (16)$$

$$q_{41}(s) = \frac{\lambda_1 \lambda_7^2}{\lambda_1 + \lambda_3 + \lambda_b} \left\{ \frac{1}{(s + \lambda_7)^2} - \frac{1}{(s + \lambda_1 + \lambda_3 + \lambda_5 + \lambda_7)^2} \right\} \quad (17)$$

$$q_{42}(s) = \frac{\lambda_3 \lambda_7^2}{\lambda_1 + \lambda_3 + \lambda_b} \left\{ \frac{1}{(s + \lambda_7)^2} - \frac{1}{(s + \lambda_1 + \lambda_3 + \lambda_5 + \lambda_7)^2} \right\} \quad (18)$$

where

$$q_{4j}(s) = \int_0^\infty e^{-st} P'_{0j} dF_7(t), \quad (j = 0, 1, 2) \quad (19)$$

Next, we define one state  $S_i$  ( $i = 0, 1, 2, 3, 4$ ) of the system at time  $t$  as that when, after the system enters this state at a time point, it will not transfer to a new state until time  $t$ . Define

that the probabilities of system state transference  $P_{0i}(t)$  ( $i = 0, 1, 2, 3, 4$ ) are the probabilities that the system starts from state  $S_0$  at time 0, enters and stays at state  $S_i$ . Their LS transforms are

$$p_{00}(s) = \frac{1 - q_{01}(s) - q_{02}(s) - q_{03}(s)}{1 - g_{00}(s)} \quad (20)$$

$$p_{01}(s) = \frac{[q_{01}(s) + q_{03}(s)q_{34}(s)q_{41}(s)][1 - q_{01}(s)]}{1 - g_{00}(s)} \quad (21)$$

$$p_{02}(s) = \frac{[q_{02}(s) + q_{03}(s)q_{34}(s)q_{42}(s)][1 - q_{02}(s)]}{1 - g_{00}(s)} \quad (22)$$

$$p_{03}(s) = \frac{q_{03}(s) - q_{03}(s)q_{34}(s)}{1 - g_{00}(s)} \quad (23)$$

$$p_{04}(s) = \frac{[q_{03}(s)q_{34}(s)][1 - q_{40}(s) - q_{41}(s) - q_{42}(s)]}{1 - g_{00}(s)} \quad (24)$$

where

$$g_{00}(s) = q_{01}(s)q_{10}(s) + q_{02}(s)q_{20}(s) + q_{03}(s)q_{34}(s)[q_{40}(s) + q_{41}(s)q_{10}(s) + q_{42}(s)q_{20}(s)] \quad (25)$$

is the recurrent first-passage time distribution of state  $S_0$ .

Recalling the steady-state probabilities of state  $S_i$  ( $i = 0, 1, 2, 3, 4$ )

$$P_i = \lim_{t \rightarrow \infty} P_{0i}(t) = \lim_{s \rightarrow 0} p_{0i}(s) \quad (26)$$

we obtain

$$P_0 = \lambda_1 + \lambda_3 + \lambda_5 \frac{1}{l_{00}} \quad (27)$$

$$P_1 = \frac{1}{\lambda_b} \left\{ \frac{\lambda_1}{\lambda_1 + \lambda_3 + \lambda_5} + \frac{\lambda_5}{\lambda_1 + \lambda_3 + \lambda_5} \frac{\lambda_1}{\lambda_1 + \lambda_3 + \lambda_b} \left[ 1 - \frac{\lambda_7^2}{(\lambda_1 + \lambda_3 + \lambda_b + \lambda_7)^2} \right] \right\} \frac{1}{l_{00}} \quad (28)$$

$$P_2 = \frac{1}{\lambda_b} \left\{ \frac{\lambda_3}{\lambda_1 + \lambda_3 + \lambda_5} + \frac{\lambda_5}{\lambda_1 + \lambda_3 + \lambda_5} \frac{\lambda_3}{\lambda_1 + \lambda_3 + \lambda_b} \left[ 1 - \frac{\lambda_7^2}{(\lambda_1 + \lambda_3 + \lambda_b + \lambda_7)^2} \right] \right\} \frac{1}{l_{00}} \quad (29)$$

$$P_3 = \left\{ \frac{1}{\lambda_b} \frac{\lambda_5}{\lambda_1 + \lambda_3 + \lambda_5} \right\} \frac{1}{l_{00}} \quad (30)$$

$$P_4 = \left\{ \frac{2}{\lambda_7} \frac{\lambda_5}{\lambda_1 + \lambda_3 + \lambda_5} \right\} \frac{1}{l_{00}} \quad (31)$$

where

$$l_{00} = \frac{1}{\lambda_b} + \frac{1}{\lambda_1 + \lambda_3 + \lambda_5} + \frac{\lambda_5}{\lambda_1 + \lambda_3 + \lambda_5} \left\{ \frac{2}{\lambda_7} + \frac{\lambda_1 + \lambda_3}{\lambda_b(\lambda_1 + \lambda_3 + \lambda_b)} \left[ 1 - \frac{\lambda_7^2}{(\lambda_1 + \lambda_3 + \lambda_b + \lambda_7)^2} \right] \right\} \quad (32)$$

is the average recurrent time at state  $S_0$ . It is evident that

$$\sum_{i=0}^4 P_i = 1. \quad (33)$$

From the Reachability Graph in Figure 2 and the state definitions, we can see that  $A1 = P_0 + P_2 + P_3 + P_4$  depicts the Availability of Robot 1 on independent operation. Similarly,  $A2 = P_0 + P_1 + P_3 + P_4$  is the Availability of Robot 2 on independent operation. The Availability of Robots 3 and 4 on cooperative motion is represented by  $P_4$ .  $P_4$  also depicts the ratio of competition between Robots 1 and 2, while the ratio of competition among Robots 1, 2 and 3 is given by  $P_0$ . Moreover, the Performability of the buffer is described by  $Pb = P_1 + P_2 + P_3$ .

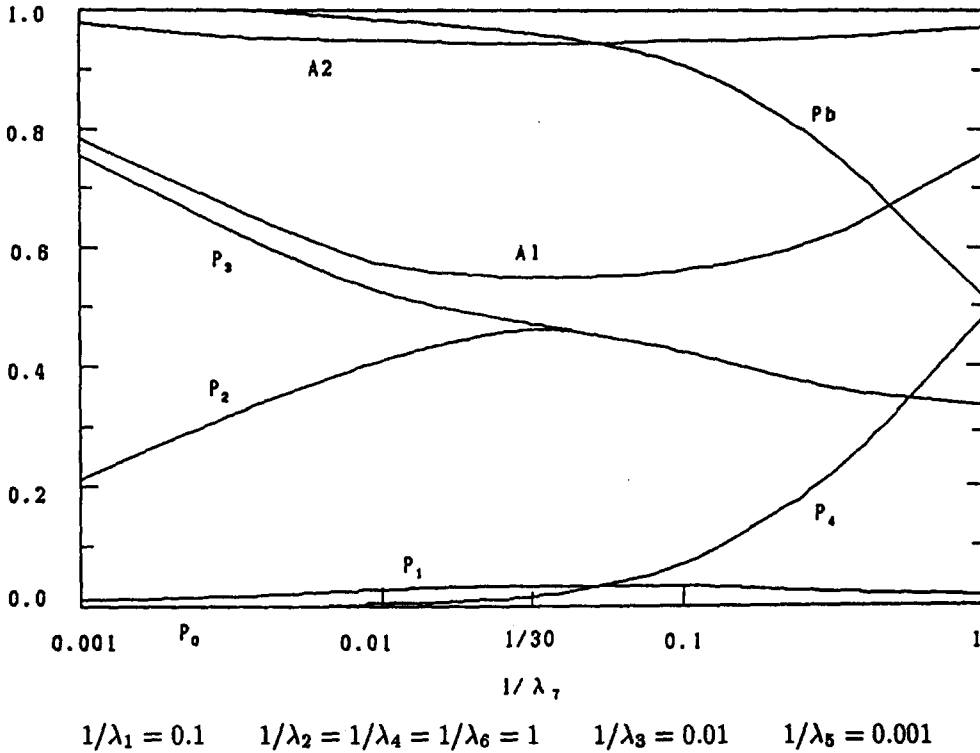


Figure 4. Numerical results of the steady-state probabilities.

## 5. NUMERICAL RESULTS

We show a numerical example when real time distributions for the firing of transitions in the ESPN model are given. Assume that the mean independent operation times of Robots 1, 2 and 3 are  $1/\lambda_1 = 0.1$  second,  $1/\lambda_3 = 0.01$  second and  $1/\lambda_5 = 0.001$  second, respectively. All the mean times of buffer operation are the same, i.e.,  $1/\lambda_2 = 1/\lambda_4 = 1/\lambda_6 = 1/\lambda_7 = 1$  second. Moreover,  $1/\lambda_7$ , the mean cooperative motion time of Robots 3 and 4, is variable from 0.001 to 1 second.

Figure 4 gives the steady-state probabilities  $P_i$  ( $i = 0, 1, 2, 3, 4$ ) and  $A1$ ,  $A2$  and  $Pb$ . It is shown from Figure 4 that the Performability of the buffer  $Pb$  decreases with  $1/\lambda_7$ . When  $1/\lambda_7$  is smaller than  $1/30$  second,  $Pb$  is very close to 1, which means that the buffer in this multi-robot system has a good performability. However, if it is more than  $1/30$  second, the Performability of the buffer  $Pb$  is worsened sharply.  $P_4$ , the Availability of Robots 3 and 4 on cooperative motion, increases with  $1/\lambda_7$ , which is consistent with the fact that the number of times to do cooperative motion is varied with the speed-up of cooperative motion. Both  $P_1$  and  $P_2$  have their maximum values at  $1/\lambda_7 = 1/30$  second, while the Availability of Robots 1 and 2 on independent operations  $A1$  and  $A2$  are shaped like a valley. This shows that a balance of independent operation rate has a great effect to the efficiency of the multi-robot system. It is interesting that  $P_2$  and  $P_3$  are almost the same when  $1/\lambda_7$  varies from  $1/30$  to 1 second. One of the most important results is that  $P_0$ , the ratio of competition among Robots 1, 2 and 3 to do buffer operation, is almost zero (at the level of  $10^{-4}$ ), which means that the competition of buffer operation does not become a problem for the system.

## 6. CONCLUSIONS

In this paper, we have developed an Extended Stochastic Petri Net model that allows firing times to non-exponential distributions, and presented an analysis method using a generalized Markov Renewal Process. We applied it to evaluate performance of a multi-robot system with parallel and cooperative motions. Our main contribution is to provide an available method to numerically solve ESPN model of concurrent systems incorporated with non-exponential distributions using a generalized MRP in which non-regeneration points may be included. The

suitability, availability and flexibility of the method for concurrent systems have been proven by the case study. The numerical results presented in this paper display the superiority over the works in [6–8]. We have also shown it is possible that, by cleverly exploiting Abstracted Partial Reachability Graph, the solution of larger model can be obtained with acceptable complexity.

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